

# Regular Languages and NFAs

CS712

Dept. of Computer Science  
Univ. of New Hampshire

Theorem Any regular language can be recognized  
by an NFA.

Proof By induction on the recursive structure  
of regular expressions.

base i)  $a \in A \quad \rightarrow \text{ } \circ \xrightarrow{a} \text{ } \circ$

ii)  $1 \quad \rightarrow \text{ } \circ \xrightarrow{1} \text{ } \circ$

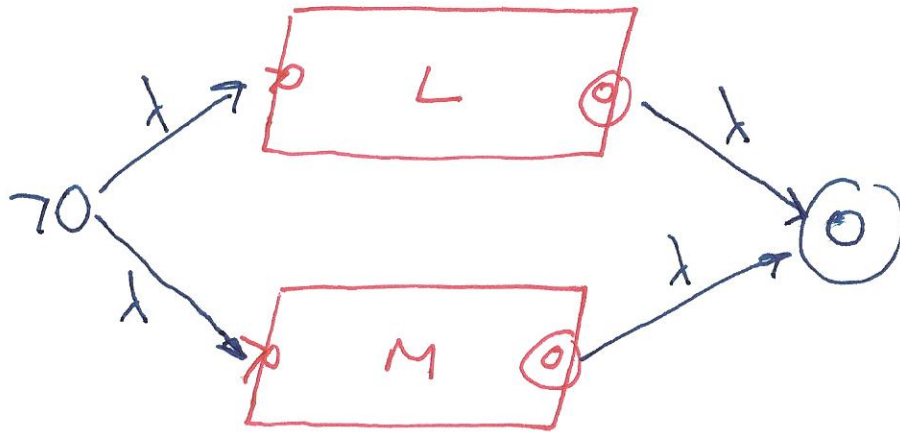
iii)  $\emptyset \quad \rightarrow \text{ } \circ$

induction step  $L$  &  $M$  are NFAs recognizing  
regular sets  $R$  &  $S$  respectively

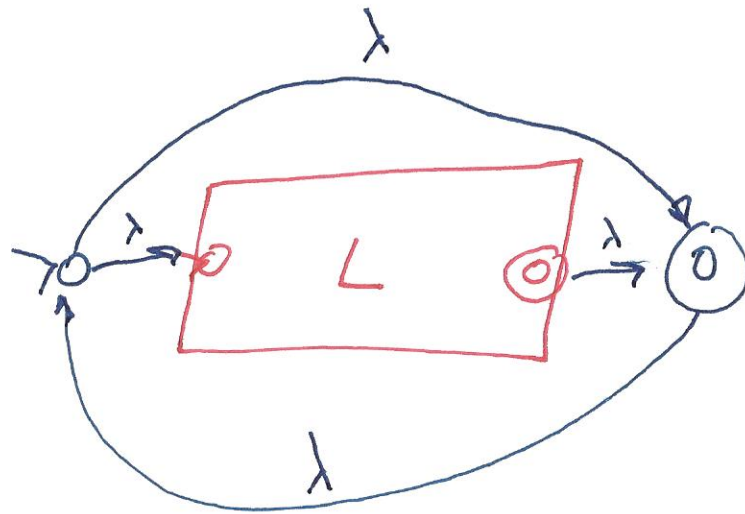
i)  $LM$



ii) L / M



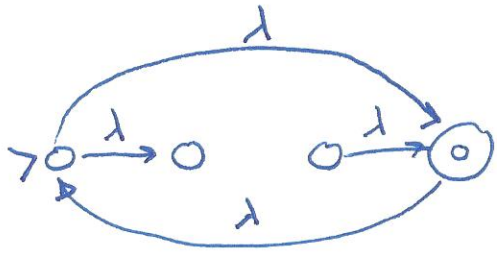
iii)  $L^*$



QED

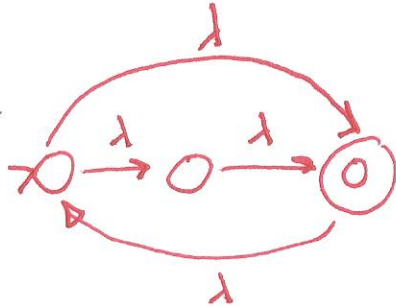
note

$\phi^*$



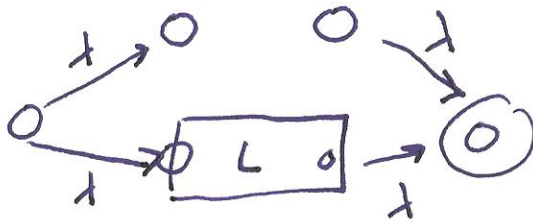
i.e.  $\phi^* = \{1\}$

$\{1\}^*$



i.e.  $\{1\}^* = \{1\}$

$\phi \cup L$



i.e.  $\phi \cup L = L$