

Signed Integer Representation

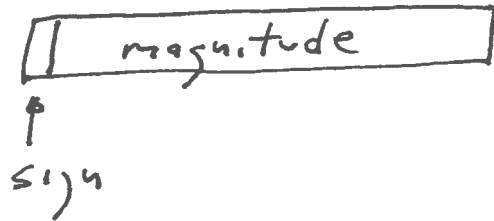
CS520

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We will consider three possibilities:

1. sign magnitude
2. one's complement
3. two's complement

sign magnitude



$$15_{10} = 0000\ 1111$$

$$-15_{10} = 1000\ 1111$$

$$\begin{array}{r} 37_{10} = 25_{16} = 00100101 \\ + (-15_{10}) \\ \hline 22_{10} \end{array}$$
$$\begin{array}{r} 00100101 \\ + 10001111 \\ \hline 00010110 = 22_{10} \end{array}$$

negate: easy

add: test sign bits

if same sign,
add magnitudes
& keep sign

if different signs,
subtract smaller
magnitude from
larger & keep
sign of larger

sign magnitude: summary

negate: easy

addition: complicated

two representations of zero

0000	0000	+0
1000	0000	-0

one's complement

express absolute value

if negative, complement all bits

↳ i.e. flip/reverse

negate: complement all bits

addition: like unsigned except if 1 carried out of leftmost bit position, add 1 to rightmost bit position

$$15_{10} = 0000\ 1111$$
$$-15_{10} = 1111\ 0000$$

$$\begin{array}{r} 37_{10} = 25_{16} \\ + (-15_{10}) \\ \hline 22_{10} \end{array}$$
$$\begin{array}{r} 0010\ 0101 \\ + 1111\ 0000 \\ \hline 0001\ 0101 \\ + 1 \\ \hline 0001\ 0110 = 16_{16} = 22_{10} \end{array}$$

$$\begin{array}{r} 10_{10} \\ + (-5_{10}) \\ \hline 5_{10} \end{array}$$
$$\begin{array}{r} 0000\ 0101 \\ + 5 \\ \hline 0000\ 0100 \\ + 1 \\ \hline 0000\ 0101 = 5 \end{array}$$

$$\begin{array}{r}
 5_{10} \\
 + (-10_{10}) \\
 \hline
 -5_{10}
 \end{array}$$

$$\begin{array}{r}
 00001010 \\
 \hline
 +10
 \end{array}$$

$$\begin{array}{r}
 0000\overset{1}{0}\overset{1}{0}1 \\
 + 1111\overset{0}{0}\overset{1}{0}1 \\
 \hline
 \end{array}$$

$$\begin{array}{r}
 11111010 = -5
 \end{array}$$

$$00000101 = +5$$

no carry, so

no second
addition

One's complement: Summary

negate: easy, do all bits at same time

addition: adding in carry could take as much time as original add

$$\begin{array}{r} 102_{10} = 66_{16} \\ + (-38_{10}) = -26_{16} \\ \hline 64_{10} \end{array}$$

$$\begin{array}{r} 00100110 \\ \hline +26 \end{array} \quad \begin{array}{r} 01100110 \\ +11011001 \\ \hline 00111111 \\ +1 \\ \hline 01000000 \end{array}$$

$$40_{16} = 64_{10}$$

one's complement: more summary

two representations for zero

0000	0000	+0
1111	1111	-0

sign bit is still leftmost bit

0	000	1111
1	111	0000

1000	0000
0111	1111

if positive,
it is too big
to fit in
8 bits
for signed
integer

two's complement

express absolute value

if negative, complement
and add 1
^
always

$$\begin{array}{r} +15_{10} = 0000\ 1111 \\ -15_{10} = 1111\ 0000 \\ \quad \quad \quad +1 \\ \hline 1111\ 0001 \end{array}$$

$$\begin{array}{r} 37_{10} = 25_{16} \\ +(-15_{10}) \\ \hline 22_{10} \end{array}$$

negate: complement and
add 1

add: same as unsigned,
ignore carry

ignored!

$$\begin{array}{r} \textcircled{1} \quad 11 \\ 0010\ 0101 \\ + 1111\ 0001 \\ \hline 0001\ 0110 \end{array}$$

$$16_{16} = 22_{10}$$

two's complement: more summary

leftmost bit still indicates sign

~~0000 0000~~

1111 1111 = -1

0000 0000

1

0000 0001 +1

all modern machines are two's complement



Moving values between different sized containers
 ^
 two's complement

promotion (small \rightarrow large): sign extend

$$\begin{array}{l}
 \underline{0000} \ \underline{1111} \rightarrow 0000 \ 0000 \ 0000 \ 1111 \\
 \underline{1111} \ \underline{0001} \rightarrow 1111 \ 1111 \ 1111 \ 0001 = -15_{10} \\
 \phantom{\underline{1111} \ \underline{0001} \rightarrow} 0000 \ 0000 \ 0000 \ 1110 \\
 \phantom{\underline{1111} \ \underline{0001} \rightarrow} + 1 \\
 \phantom{\underline{1111} \ \underline{0001} \rightarrow} 0000 \ 0000 \ 0000 \ 1111 = +15_{10}
 \end{array}$$

demotion (large \rightarrow small): chop \rightarrow overflow?

$$\begin{array}{r}
 \cancel{0000} \ \cancel{0000} \ 0000 \ 1111 \\
 \cancel{1111} \ \cancel{1111} \ 1111 \ 0001 \\
 \phantom{\cancel{0000} \ \cancel{0000} \ } 0000 \ 1110 \\
 \phantom{\cancel{0000} \ \cancel{0000} \ } + 1 \\
 \hline
 0000 \ 1111
 \end{array}$$

$$\begin{array}{r}
 257_{10} \ \cancel{0000} \ \cancel{0001} \ 0000 \ 0001 = 257_{10} \\
 \phantom{257_{10} \ } \phantom{\cancel{0000} \ \cancel{0001} \ } 0000 \ 0001 = 1_{10}
 \end{array}$$

overflow if the chopped bits are not all the same.

$$\begin{array}{r}
 -257_{10} \ 1111 \ 1110 \ 1111 \ 1110 \\
 \phantom{-257_{10} \ } + 1 \\
 \hline
 \cancel{1111} \ \cancel{1110} \ 1111 \ 1111 = -257_{10} \\
 \phantom{\cancel{1111} \ \cancel{1110} \ } 1111 \ 1111 = -1_{10}
 \end{array}$$