CS725/8258 [T725] Lecture 4 Networking Fundamentals

September 9, 2024

Anatomy of a router/switch







Header info (destination address)

Store and Forward

- Intermediate nodes receive, store, and forward packets
 - storing and retransmission adds fixed delay
- Output conflict multiple packets waiting to be forwarded on the same output
 - queues/buffers used to store packets waiting transmission
 - queuing adds variable delay
 - potential for packet loss due to buffer overflow

A bit of history...

- Local area networks (late 80's, early 90's)
 - (then) based on broadcast & select medium



- (today) a network of Ethernet (L2) switches



Broadcast & select medium

Comparison

Routed networks



- topology driven by geography
- long distances (high latency)
- need for scalability
- location-related addresses
- routing



Broadcast & select



- everyone connected to everyone
- short distances (low latency)
- lesser need for scalability
- arbitrary addresses
- address discovery





Internet - a network to INTERconnect NETworks



Today, L2 switching and L3 routing are often mixed together in devices capable of processing multiple layers



Networking Basics

- Goals: IP/MAC addresses, to basics of routing/switching
- Assumptions:
 - packet switched network
 - nodes attached to a L2
 broadcast-and-select network
 - each node "has" a 6-byte <u>MAC</u> and a 4-byte <u>IP</u> addresses

Goals: IP/MAC addresses, fundamental protocols, subnetting,

<u>C</u> and



Address Resolution

A has a packet with IP destination address B, A needs B's MAC address to deliver the packet

Solution: ARP - Address Resolution Protocol

Problem: Find MAC address of a node with a given IP address



ARP

A has a packet with IP destination address B, A needs B's MAC address to deliver the packet





ARP

Step 1: "Who has ..." broadcasted to everyone by A



Who has IP = B?





ARP request processed in B (MAC b-cast, ARP IP match)





Step 2: B sends L2 unicast response to A (MAC a) Due to the broadcast nature of the underlying medium, the response will be seen by all other nodes B is at MAC b but dropped (MAC mismatch) (L2 unicast to MAC a) Broadcast and Select medium Not Not for for me! me! IP A В MAC a b С





Step 3: A sends data using L2 unicast to MAC b





ARP request processed in B (MAC b-cast, ARP IP match)

ARP request processed in C (MAC b-cast, IP mismatch)

ARP response processed in C (MAC unicast, MAC mismatch)

IP packet processed in B (MAC unicast, MAC match,



Subnetting

- Problem: deciding whether an IP address belongs to a specific subset of IP addresses
- Solution: nodes on a subnet (and only those) have IP addresses within a specific range

an IP address t of IP addresses et (and only those) specific range





Subnetting

- Problem: deciding whether an IP address belongs to a specific subset of IP addresses
- Solution: nodes on a subnet (and only those) have IP addresses within a specific range
- Simplifying HW implementation: subnets (ranges of IP addresses that can be placed only in a specific, constrained way):

Subnet ID

Prefix length (bits)







IP Pretix

- specifies a range of consecutive IP addresses
- consists of a subnet id and a length (len) e.g., 132.177.4.0/22
- An IP address belongs to a range specified by a prefix if its first len bits are equal to those in the subnet id $(s_i = b_i \text{ for all } i \leq len):$





prefix

Example questions

Range of 132.177.4.0/24? 132.177.4.0 → 132.177.4.255

Prefix for range $132.177.0.0 \rightarrow 132.177.255.255?$ 132.177.0.0/16

Example questions

Range of 132.177.4.0/26? 132.177.4.0 → 132.177.4.63

Prefix for range $132.177.2.192 \rightarrow 132.177.2.223?$ 132.177.2.192/27

Combining prefixes

132.177.0.0/24 132.177.1.0/24

Range 132.177.0.0 - 132.177.0.255 Range 132.177.1.0 - 132.177.1.255

Combining prefixes

 132.177.0.0/24
 132.177.0.0/23
 Range 132.177.0.0 - 132.177.1.255

 132.177.1.0/24
 132.177.0.0/23
 Range 132.177.0.0 - 132.177.1.255



Combining prefixes

132.177.0.0/24 132.177.0.0/23 132.177.1.0/24 132.177.2.0/24 132.177.2.0/23 132.177.3.0/24 132.177.4.0/24 132.177.4.0/23 132.177.5.0/24 132.177.6.0/24 132.177.6.0/23 132.177.7.0/24 132.177.8.0/24 132.177.8.0/23 132.177.9.0/24 132.177.10.0/24 132.177.10.0/23 132.177.11.0/24 132.177.12.0/24 132.177.12.0/23 132.177.13.0/24 132.177.14.0/24 132.177.14.0/23 132.177.15.0/24



Subnet Mask

- Another way to specify prefix length



A 32 bit, IP address-like value whose binary representation has binary ones in in bits corresponding to the subnet id bits.

Subnet Mask Representation

Typically represented using decimal dotted notation: 255.255.0.0

Subnet length and subnet mask (netmask) are for all practical purposes equivalent:

> 132.177.4.0 with netmask 255.255.255.0 is equivalent to 132.177.4.0 /24

Example questions

- Subnet mask for prefix length /17?
 17 ones
 - 11111111 111111
 - 255 255
- Prefix length for subnet mask 255.255.255.192?
 26 ones

11111111 1111111

255 255





Domain Name Service



Domain

Top Level Domain (TLD) WWW.CS.unh.edu

Hostname

Domain Name Service

- Mapping between hostnames and IP addresses:
 - one-to-one, one-to-many, many-to-one, or many-to-many?
 - mapping in both directions
- Possible solutions:
 - centralized database
 - fully distributed database